Optimal Low-Rank Matrix Completion: Semidefinite Relaxations and Eigenvalue Disjunctions

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(2)

Problem Setting

Find low-rank X that approx. fits observed entries $\mathcal I$ in noisy data A:

$$\min_{\boldsymbol{X} \in \mathbb{R}^{n \times m}} \frac{1}{2} \sum_{(i,j) \in \mathcal{T}} (X_{i,j} - A_{i,j})^2 + \frac{1}{2\gamma} \|\boldsymbol{X}\|_F^2 \quad \text{s.t. Rank} (\boldsymbol{X}) \le k \quad (1)$$

 $\underset{\boldsymbol{X} \in \mathbb{R}^{n \times m}}{\min} \|\boldsymbol{X}\|_{F}^{2} \quad \text{s.t. } X_{i,j} = A_{i,j} \forall (i,j) \in \mathcal{I}, \text{ Rank} (\boldsymbol{X}) \leq k$

Summary of Contributions

- 1. Branch-and-bound by eigenvector disjunctions
- 2. Tighter semidefinite inequalities via determinant minors
- 3. Restricted **alternating minimization** at each node
- ⇒ Solves (nonconvex) low-rank problem to **global optimality**

Matrix Perspective Reformulation Technique

Reformulate (1) exactly as [c.f. 2]: $\min_{\mathbf{V} \subset \mathcal{V}^k} \min_{\mathbf{V} \subset \mathbb{D}^n \times m} \frac{1}{2} \sum (X_{i,j} - A_{i,j})^2 + \frac{1}{2\gamma} \operatorname{tr}(\boldsymbol{\Theta})$

3. Alternating Minimization

Following [3], express X = UV, iterate until local minima reached

- Improvement: compute high-quality solutions at child nodes [e.g., in 4]
- Restrict U in feasible region of child node: linear constraints, and SOC approximation of $U^{\top}U \preceq \mathbb{I}$ [c.f. 1]
- *Randomly* initiate altmin, with probability decaying with tree depth.

- ▶ Isolates nonconvexity of Rank $(\mathbf{X}) \leq k$ in set of rank-k projections \mathcal{Y}_n^k
- ▶ Relaxing to $Y \in Conv(\mathcal{Y}_n^k) = \{ \mathbf{0} \leq Y \leq \mathbb{I}, tr(Y) \leq k \}$ gives a SDP!

1. Eigenvector Disjunctions

Introduce
$$\boldsymbol{U} \in \mathbb{R}^{n \times k}$$
, and relax $\boldsymbol{Y} = \boldsymbol{U}\boldsymbol{U}^{\top}$ to $\boldsymbol{Y} \succeq \boldsymbol{U}\boldsymbol{U}^{\top}$:

$$\min_{\substack{\boldsymbol{Y} \in \mathsf{Conv}(\mathcal{Y}_{n}^{k}) \\ \boldsymbol{U} \in \mathbb{R}^{n \times k}}} \min_{\substack{\boldsymbol{X} \in \mathbb{R}^{n \times m} \\ \boldsymbol{\Theta} \in \mathcal{S}^{m}}} \frac{1}{2} \sum_{(i,j) \in \mathcal{I}} (X_{i,j} - A_{i,j})^{2} + \frac{1}{2\gamma} \mathsf{tr}\left(\boldsymbol{\Theta}\right)$$

$$\mathbf{S.t.} \quad \begin{pmatrix} \boldsymbol{Y} & \boldsymbol{X} \\ \boldsymbol{X}^{\top} & \boldsymbol{\Theta} \end{pmatrix} \succeq \boldsymbol{0}, \ \boldsymbol{Y} \succeq \boldsymbol{U}\boldsymbol{U}^{\top}$$

$$(3)$$

► If solution (\hat{Y}, \hat{U}) has $\hat{Y} \preceq \hat{U}\hat{U}^{\top}$, done!

Else, there is eigenvector
$$\boldsymbol{x}$$
 such that $\boldsymbol{x}^{\top} (\boldsymbol{Y} - \boldsymbol{U}\boldsymbol{U}^{\top}) \boldsymbol{x} > 0$. Cut:
 $\boldsymbol{x}^{\top} \boldsymbol{Y} \boldsymbol{x} \leq \boldsymbol{x}^{\top} \boldsymbol{U} \boldsymbol{U}^{\top} \boldsymbol{x} = \| \boldsymbol{U}^{\top} \boldsymbol{x} \|_{2}^{2}$ (4)

▶ k = 1: Following [6], linearize $u \mapsto u^2$ on [-1, 1] as below:

$$P \leq \begin{cases} f_L(u;u_0) := -u + uu_0 + u_0 & \text{if } u \in [-1,u_0] \\ f_R(u;u_0) := u + uu_0 - u_0 & \text{if } u \in (u_0,1] \end{cases}$$

and set $u_0 = \hat{\boldsymbol{U}}^\top \boldsymbol{x}, u = \boldsymbol{U}^\top \boldsymbol{x}$.

 u^2

► k > 1: Yields 2^k disjunctive regions:

$$\bigvee_{L\subseteq[k]} \left\{ \left(\boldsymbol{U},\boldsymbol{Y}\right) \middle| \begin{array}{l} \boldsymbol{U}_{j}^{\top}\boldsymbol{x} \in [-1, \hat{\boldsymbol{U}}_{j}^{\top}\boldsymbol{x}] & \forall j \in L, \\ \boldsymbol{U}_{j}^{\top}\boldsymbol{x} \in (\hat{\boldsymbol{U}}_{j}^{\top}\boldsymbol{x}, 1] & \forall j \in [k] \setminus L, \\ \boldsymbol{x}^{\top}\boldsymbol{Y}\boldsymbol{x} \leq \sum_{j \in L} f_{L}(\boldsymbol{U}_{j}^{\top}\boldsymbol{x}; \hat{\boldsymbol{U}}_{j}^{\top}\boldsymbol{x}) & + \sum_{j \in [k] \setminus L} f_{R}(\boldsymbol{U}_{j}^{\top}\boldsymbol{x}; \hat{\boldsymbol{U}}_{j}^{\top}\boldsymbol{x}) \end{array} \right\} \quad (5)$$

Note: different approximations \rightarrow different branching factor!

Computational Experiments: Root Node Relaxations

Imposing some Shor SDP constraints provide tradeoff between relaxation strength and solve time!



Figure: Root node relative gap (left) and time taken (right) at the root node for rank-1 matrix completion with $2kn \log_{10}(n)$ filled entries, with low regularization ($\gamma = 80.0$).

Computational Experiments: Scalability and Accuracy

► As *n* increases, SDP relaxations are tighter but take longer to solve!



Figure: Relative optimality gap for rank-1 (left) and rank-2 (right) matrix completion



2. Tight SDP Inequalities via Determinant Minors

In rank-1 matrices, each 2-by-2 minor has determinant 0. Taking the moment matrix of $\begin{pmatrix} 1 & X_{i_1,j_1} & X_{i_1,j_2} & X_{i_2,j_1} & X_{i_2,j_2} \end{pmatrix}$ gives:

$$\det \begin{pmatrix} X_{i_1,j_1} & X_{i_1,j_2} \\ X_{i_2,j_1} & X_{i_2,j_2} \end{pmatrix} = 0 \implies \begin{pmatrix} 1 & X_{i_1,j_1} & X_{i_1,j_2} & X_{i_2,j_1} & X_{i_2,j_2} \\ X_{i_1,j_1} & W_{i_1,j_1} & V_{i_1,(j_1,j_2)}^1 & V_{(i_1,i_2),j_1}^2 & V_{(i_1,i_2),(j_1,j_2)}^3 \\ X_{i_1,j_2} & V_{i_1,(j_1,j_2)}^1 & W_{i_1,j_2} & V_{(i_1,i_2),(j_1,j_2)}^3 & V_{(i_1,i_2),j_2}^2 \\ X_{i_2,j_1} & V_{(i_1,i_2),j_1}^2 & V_{(i_1,i_2),(j_1,j_2)}^3 & W_{i_2,j_1} & V_{i_2,(j_1,j_2)}^1 \\ X_{i_2,j_2} & V_{(i_1,i_2),(j_1,j_2)}^3 & V_{i_2,(j_1,j_2)}^2 & V_{i_2,(j_1,j_2)}^1 & W_{i_2,j_2} \end{pmatrix} \succeq \mathbf{C}$$

- ▶ Let $W_{i,j}$ model $X_{i,j}^2$ in objective and constraints;
- Link $\Theta_{i,j}$ (which models $\sum_l X_{i,l} X_{j,l}$) to appropriate W and V terms:

$$\Theta_{j_1,j_2} = \sum_{i \in [n]} V_{i,(j_1,j_2)}^1, \quad \forall \ j_1 < j_2 \in [m], \qquad \Theta_{j,j} = \sum_{i \in [n]} W_{i,j}, \quad \forall \ j \in [m].$$

- Extend to rank k > 1 by modelling X as sum of rank-1 slices $\sum_{t \in [k]} X^t$:
 - Requires Shor relaxation on the moment matrix of $\begin{pmatrix} 1 & X_{i,j}^1 & X_{i,j}^2 & \dots & X_{i,j}^k \end{pmatrix}$
 - New terms $H_{i,j}^{t_1,t_2}$ modelling $X_{i,j}^{t_1}X_{i,j}^{t_2}$
- **Key insight:** modelling *some* minors \rightarrow cutting planes for (3)

with $pkn \log_{10}(n)$ filled entries, with low regularization ($\gamma = 80.0$).

Outperforms B-M at root node and MFSGD [5] in terms of MSE



Figure: (Absolute) MSE improvement against alternating minimization at the the root node (left) and MatrixFactorization() (right) for rank-*k* matrix completion with $2kn \log_{10}(n)$ filled entries, with low regularization ($\gamma = 80.0$).

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